

## Verifying Safety Properties in Assembly Code without Compiler Support

CENTER FOR HYBRID & EMBEDDED SOFTWARE SYSTEMS

<http://chess.eecs.berkeley.edu>

### Problem:

### We need to ensure binary code is safe to execute

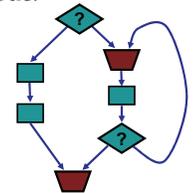
- Is it memory-safe/type-safe?
- Does it use APIs properly?
- Does it obey resource constraints?
- What about other security properties?
- Many software vendors won't release source code, but binary code is hard to analyze.
- We have security tools that analyze source code. Can we use these to help us prove that binary code is safe?



Idea: Proof-Carrying Code that extends easily to different source-code analyses.

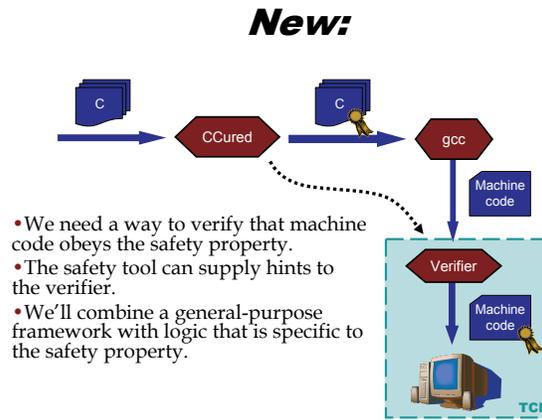
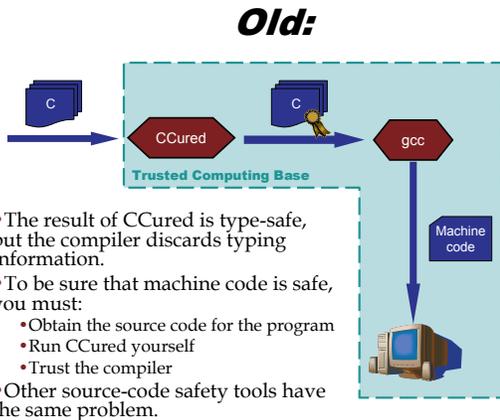
### How can we do this?

1. We need a dependently-typed assembly language to express the complex invariants of the source-level analysis.
  - Take advantage of off-the-shelf compilers.
  - Strategy: Abstract interpretation over the assembly code.
    - Track the type of each register.
    - Challenge: **Join points**.
      - How do we deal with dependent values?
      - What information can we throw away to improve performance?
2. We need type inference so that we don't need a type-preserving compiler.



### Prototype Implementation

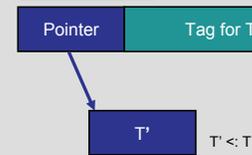
- We can verify most of CCured's output.
- In the spec95 "go" benchmark, it takes about 0.40s to verify each function.
- CCured supplies us with a few hints (e.g. function signatures.)



### CCured

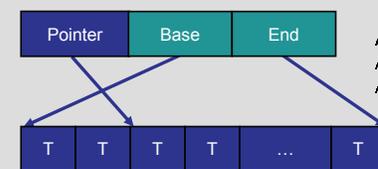
- A source-to-source translator that makes C programs type-safe.
- Inserts runtime checks before unsafe operations.
- Adds *metadata* to some pointers to support these checks.

RTTI pointer to T (for checked downcasts)



- These "fat pointers" are dependent types!
- We must track dependencies between registers.
- Challenge: Changes to memory structures.

Sequence pointer to T (for array bounds checks)



$$\begin{aligned}
 & (base \leq end) \\
 & \wedge (end - base) \bmod \text{sizeof}(T) = 0 \\
 & \wedge (pointer - base) \bmod \text{sizeof}(T) = 0 \\
 & \wedge \forall i. (base \leq (pointer + i \cdot \text{sizeof}(T)) < end) \Rightarrow \\
 & \quad ((pointer + i \cdot \text{sizeof}(T)) \text{ inPtr } T)
 \end{aligned}$$

### Future Work

1. Extend the implementation to more safety tools.
  - Type safety for OO languages.
  - Security properties using Cqual.
2. Generate low-level proofs of safety properties.
  - Foundational PCC
3. Stronger dependent types
  - Allow dependencies among different memory objects.
  - Immutable dependencies.