

Introduction to Embedded Systems

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EECS 124
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Lecture 12: Temporal Logic

Specification, Verification, and Control

Specification

A mathematical statement of the goal to be achieved

Verification

Does the system achieve its goal, as designed?

Controller Synthesis

Given an incomplete design, synthesize a strategy to complete the system so that it achieves its goal

Specification, Verification, and Control

Specification

TEMPORAL LOGIC

A mathematical statement of the goal to be achieved

Verification

REACHABILITY ANALYSIS

Does the system achieve its goal, as designed?

Controller Synthesis

Given an incomplete design, synthesize a strategy to complete the system so that it achieves its goal

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Temporal Logic

- A mathematical way to express properties of a system over time
 - E.g., Behavior of an FSM or Hybrid System
- Many flavors of temporal logic
 - Propositional temporal logic
 - Real-time temporal logic
- Amir Pnueli won ACM Turing Award, in part, for the idea of using temporal logic for specification

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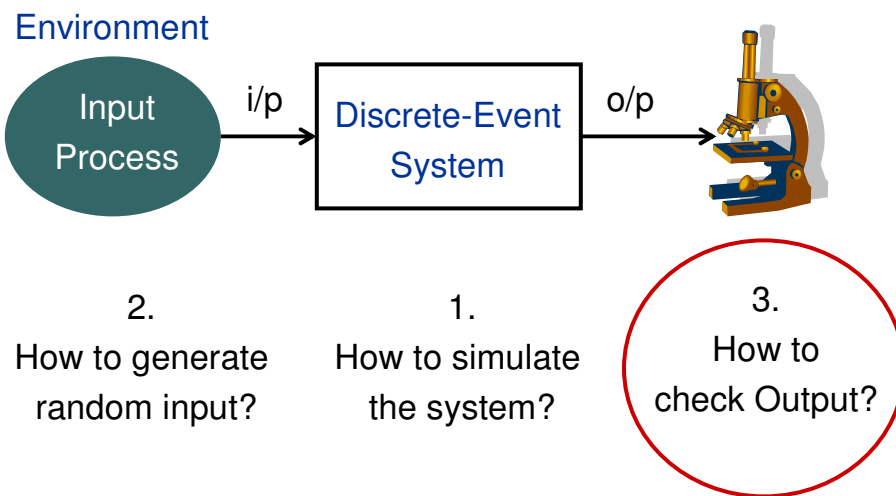
Example: Specification of the *SpaceWire* Protocol (European Space Agency standard)

8.5.2.2 ErrorReset

- a. The *ErrorReset* state shall be entered after a system reset, after link operation is terminated for any reason or if there is an error during link initialization.
- b. In the *ErrorReset* state the Transmitter and Receiver shall all be reset.
- c. When the reset signal is de-asserted the *ErrorReset* state shall be left unconditionally after a delay of 6,4 μ s (nominal) and the state machine shall move to the *ErrorWait* state.
- d. Whenever the reset signal is asserted the state machine shall move immediately to the *ErrorReset* state and remain there until the reset signal is de-asserted.

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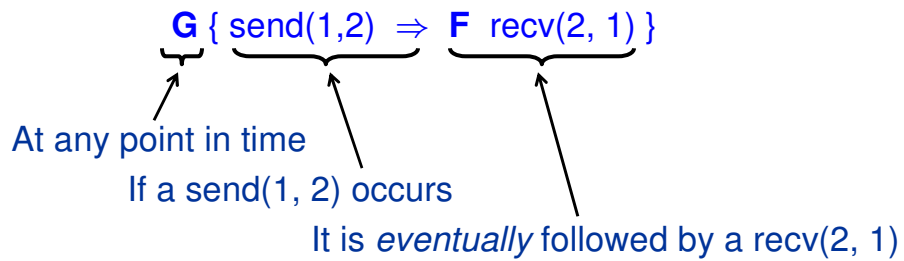
Recap: Simulating a Discrete-Event System



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Propositional Temporal Logic

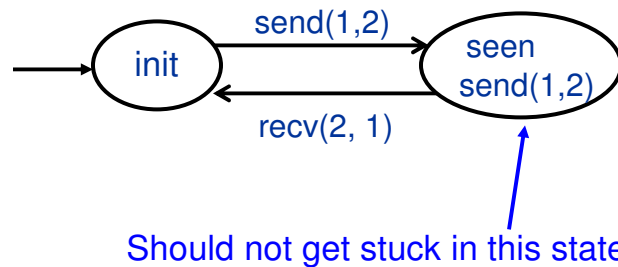
Every $\text{send}(1, 2)$ is eventually followed by a $\text{recv}(2, 1)$



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Propositional Temporal Logic

Every $\text{send}(1, 2)$ is eventually followed by a $\text{recv}(2, 1)$

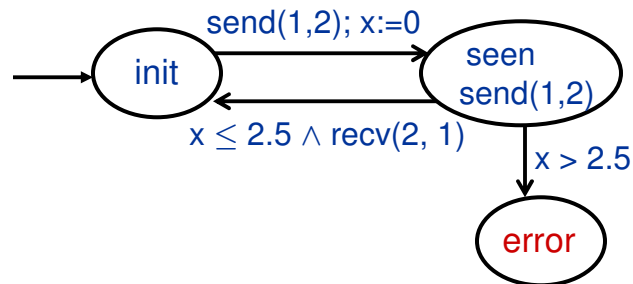
$$\mathbf{G} \{ \text{send}(1,2) \Rightarrow \mathbf{F} \text{recv}(2, 1) \}$$


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Real-Time Temporal Logic

Every $\text{send}(1, 2)$ is followed by a $\text{recv}(2, 1)$ within 2.5 ms

$$\mathbf{G} \{ \text{send}(1,2) \Rightarrow \mathbf{F}_{\leq 2.5} \text{recv}(2, 1) \}$$



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Property on a Single State

A Boolean expression over system state or input/output symbols

- $\text{send}(1,2)$: packet sent from node 1 to node 2
- $x > 2.5$: clock variable x exceeds 2.5

We will visualize each such Boolean expression by

- a distinct color
- using a name such as p , q , ...



$\text{send}(1,2)$



$x > 2.5$

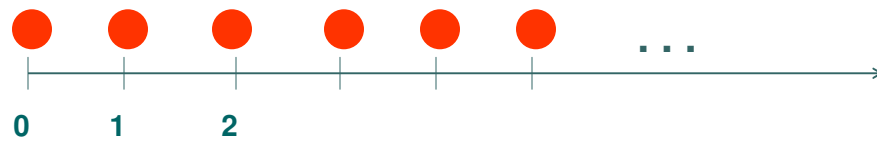
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Globally p : $G p$

$G p$ is true in a state if p holds at all points of time (along the path) starting from that state

$p = \bullet$

$G p$ holds in the initial state iff



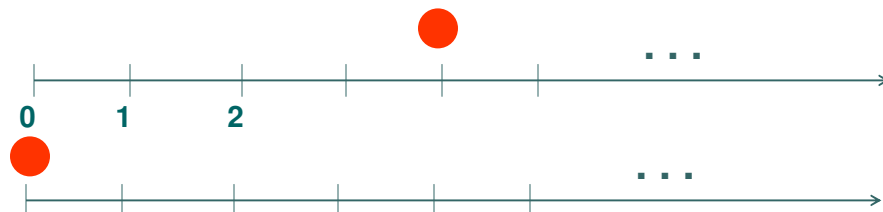
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Eventually p : $F p$

$F p$ is true in a state if p holds at some point in the future, starting from that state

$p = \bullet$

If $F p$ holds in the initial state, we could have the scenario:



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Next p: X p

X p is true along a path starting in a state if p holds in the next state

p = ●

Suppose X p holds in state at t = 2



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Nesting of Formulas

p need not be just a Boolean formula.

It can be a temporal logic formula itself!

p = ●

“X p holds in all states, starting from initial state”

How can we write this in temporal logic?

How do we draw this?

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Alternate Notation

Sometimes you'll see alternative notation in the literature:

G □

F ◇

X ○

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Examples: What do they mean?

○ $G F p$

○ $F G p$

○ $G(p \rightarrow F q)$

○ $F(p \rightarrow (X X q))$

Remember:

Gp p holds in all states

Fp p holds eventually

Xp p holds in the next state

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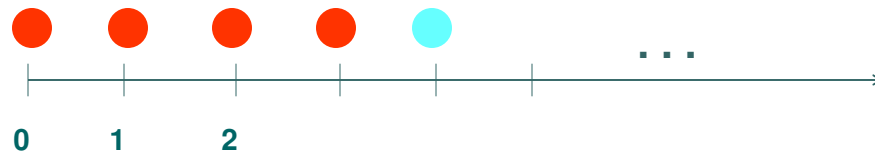
p Until q : $p \text{ U } q$

$p \text{ U } q$ is true in a state if

- q is true in some state reachable from s
- p is true in all states from s until q holds

$p =$ ● $q =$ ●

Suppose $p \text{ U } q$ holds in initial state



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Temporal Operators & Relationships

G , F , X , U : All express properties along system traces

- Can you express $G p$ purely in terms of F , p , and Boolean operators ?
- How about F in terms of U ?
- What about X in terms of G , F , or U ?

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Examples: Write in Temporal Logic

1. “Whenever the iRobot is at the ramp-edge (cliff), eventually it moves 5 cm away from the cliff.”
 - p – iRobot is at the cliff
 - q – iRobot is 5 cm away from the cliff
2. “Whenever the distance between cars is less than 2m, cruise control is deactivated”
 - p – distance between cars is less than 2 m
 - q – cruise control is active

Ex. write the SpaceWire specs. in Temporal Logic